Numerical methods

System of linear equations (cont.)

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- 1. Determinants and Cramer's rule
- 2. Iterative method for solving linear systems
- 3. Convergence, iteration matrix, stop criteria
- 4. Jacobi method, Gauss-Seidel method
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Determinant of the second order

$$\mathbf{a}_{11}\mathbf{x}_1 + \mathbf{a}_{12}\mathbf{x}_2 = \mathbf{b}_1$$

 $\mathbf{a}_{21}\mathbf{x}_1 + \mathbf{a}_{22}\mathbf{x}_2 = \mathbf{b}_2$

$$(a_{11}a_{22} - a_{12}a_{21})x_1 = b_1a_{22} - b_2a_{12}$$

$$(a_{11}a_{22} - a_{12}a_{21})x_2 = a_{11}b_2 - a_{21}b_1$$

$$\mathbf{x}_1 = \frac{\mathbf{b}_1 \mathbf{a}_{22} - \mathbf{b}_2 \mathbf{a}_{12}}{\mathbf{a}_{11} \mathbf{a}_{22} - \mathbf{a}_{12} \mathbf{a}_{21}}, \quad \mathbf{x}_2 = \frac{\mathbf{a}_{11} \mathbf{b}_2 - \mathbf{a}_{21} \mathbf{b}_1}{\mathbf{a}_{11} \mathbf{a}_{22} - \mathbf{a}_{12} \mathbf{a}_{21}}.$$

Definition: Let

$$\mathbf{A} = \begin{pmatrix} \mathbf{a}_{11} & \mathbf{a}_{12} \\ \mathbf{a}_{21} & \mathbf{a}_{22} \end{pmatrix}$$

be a matrix over a field F.

Then the term $a_{11}a_{22} - a_{12}a_{21}$ is called **determinant** (of second order)

of matrix A

and we denote it as $a_{11}a_{22} - a_{12}a_{21} = \begin{vmatrix} a_{11} & a_{12} \\ a_{21} & a_{22} \end{vmatrix} = |A|.$

$$\mathbf{x}_{1} = \frac{\begin{vmatrix} \mathbf{b}_{1} & \mathbf{a}_{12} \\ \mathbf{b}_{2} & \mathbf{a}_{22} \end{vmatrix}}{\begin{vmatrix} \mathbf{a}_{11} & \mathbf{a}_{12} \\ \mathbf{a}_{21} & \mathbf{a}_{22} \end{vmatrix}} \qquad \mathbf{x}_{2} = \frac{\begin{vmatrix} \mathbf{a}_{11} & \mathbf{b}_{1} \\ \mathbf{a}_{21} & \mathbf{b}_{2} \end{vmatrix}}{\begin{vmatrix} \mathbf{a}_{11} & \mathbf{a}_{12} \\ \mathbf{a}_{21} & \mathbf{a}_{22} \end{vmatrix}}$$

Determinant of the third order

$$a_{11}x_1 + a_{12}x_2 + a_{13}x_3 = b_1$$

$$a_{21}x_1 + a_{22}x_2 + a_{23}x_3 = b_2$$

$$a_{31}x_1 + a_{32}x_2 + a_{33}x_3 = b_3$$

Lets multiply the first equation by $a_{22}a_{33} - a_{23}a_{32}$, the second equation by $a_{13}a_{32} - a_{12}a_{33}$, the third equation by $a_{12}a_{23} - a_{13}a_{22}$, and then sum-up the first two with the third one and we get:

$$(a_{11}a_{22}a_{33} - a_{11}a_{23}a_{32} + a_{13}a_{21}a_{32} - a_{12}a_{21}a_{33} + a_{12}a_{23}a_{31} - a_{13}a_{22}a_{31}) x_1$$

$$= b_1a_{22}a_{33} - b_1a_{23}a_{32} + b_2a_{13}a_{32}$$

 $-b_2a_{12}a_{33}+b_3a_{12}a_{23}-b_3a_{13}a_{22}$

Definition: Let A is a square matrix of the order of 3 with the elements from the field F. Then the **determinant** (of the third order) of matrix A is the term

 $\begin{aligned} a_{11}a_{22}a_{33} - a_{11}a_{23}a_{32} + a_{13}a_{21}a_{32} \\ - a_{12}a_{21}a_{33} + a_{12}a_{23}a_{31} - a_{13}a_{22}a_{31} = \\ \begin{vmatrix} a_{11} & a_{12} & a_{13} \\ a_{21} & a_{22} & a_{23} \\ a_{31} & a_{32} & a_{33} \end{vmatrix} = |A|. \end{aligned}$

If we denote

$$A_{1} = \begin{pmatrix} b_{1} & a_{12} & a_{13} \\ b_{2} & a_{22} & a_{23} \\ b_{3} & a_{32} & a_{33} \end{pmatrix}, A_{2} = \begin{pmatrix} a_{11} & b_{1} & a_{13} \\ a_{21} & b_{2} & a_{23} \\ a_{31} & b_{3} & a_{33} \end{pmatrix}, A_{3} = \begin{pmatrix} a_{11} & a_{12} & b_{1} \\ a_{21} & a_{22} & b_{2} \\ a_{31} & a_{32} & b_{3} \end{pmatrix},$$

then

$$x_1 = \frac{|A_1|}{|A|}, \quad x_2 = \frac{|A_2|}{|A|}, \quad x_3 = \frac{|A_3|}{|A|}.$$

Cramer's rule

a ₁₁	a ₁₂	a ₁₃
a ₂₁	a ₂₂	a ₂₃
a ₃₁	a ₃₂	a ₃₃

a ₁₁	a ₁₂	a ₁₃
a ₂₁	^a 22	a ₂₃
a ₃₁	a ₃₂	^a 33
a ₁₁	a ₁₂	a ₁₃
a ₂₁	a ₂₂	a ₂₃

^a 11	a ₁₂	a ₁₃
^a 21	^a 22	a ₂₃
a ₃₁	^a 32	^a 33
a ₁₁	a ₁₂	^a 13
a ₂₁	a ₂₂	a ₂₃

^a 11	a ₁₂	a ₁₃
^a 21	^a 22	a ₂₃
^a 31	^a 32	^a 33
a ₁₁	^a 12	^a 13
a ₂₁	a ₂₂	^a 23

 $\mathbf{a_{11}a_{22}a_{33}} - \mathbf{a_{11}a_{23}a_{32}} + \mathbf{a_{13}a_{21}a_{32}} - \mathbf{a_{12}a_{21}a_{33}} + \mathbf{a_{12}a_{23}a_{31}} - \mathbf{a_{13}a_{22}a_{31}}$

a ₁₁	a ₁₂	a ₁₃
a_{21}	a ₂₂	a ₂₃
a ₃₁	a ₃₂	a ₃₃
a ₁₁	a ₁₂	a ₁₃
a ₂₁	a ₂₂	a ₂₃

 $\mathbf{a_{11}a_{22}a_{33}} - \mathbf{a_{11}a_{23}a_{32}} + \mathbf{a_{13}a_{21}a_{32}} - \mathbf{a_{12}a_{21}a_{33}} + \mathbf{a_{12}a_{23}a_{31}} - \mathbf{a_{13}a_{22}a_{31}}$

a ₁₁	a ₁₂	^a 13
a ₂₁	^a 22	a ₂₃
^a 31	a ₃₂	a ₃₃
a ₁₁	a ₁₂	a ₁₃
a ₂₁	a ₂₂	a ₂₃

a ₁₁	a ₁₂	^a 13
a ₂₁	^a 22	^a 23
^a 31	^a 32	a ₃₃
^a 11	a ₁₂	a ₁₃
a ₂₁	a ₂₂	a ₂₃

a ₁₁	a ₁₂	^a 13
a ₂₁	^a 22	^a 23
^a 31	^a 32	^a 33
^a 11	^a 12	a ₁₃
^a 21	a ₂₂	a ₂₃

Determinant of a square matrix

If $\mathbf{A} = \begin{bmatrix} a_{ij} \end{bmatrix}$ is a square matrix of type $n \times n$, then the **determinant** is the exactly defined number we denote as $|\mathbf{A}|$.

The minor A_{ij} is defined to be the determinant of the $(n-1)\times(n-1)$ matrix that results from **A** by removing the *i*th row and the *j*th column.

The determinant itself is recursively defined as follows: If n = 1, then the determinant of matrix $\mathbf{A} = \begin{bmatrix} a_{11} \end{bmatrix}$ of type 1×1 is simply $|\mathbf{A}| = a_{11}$. If $n \ge 2$, then for each row index *i* it holds: $|\mathbf{A}| = (-1)^{i+1} a_{i1} |\mathbf{A}_{i1}| + (-1)^{i+2} a_{i2} |\mathbf{A}_{i2}| + \dots + (-1)^{i+n} a_{in} |\mathbf{A}_{in}|$.

This is called

the Laplace expansion along the ith row.

If we have a LU decomposition of matrix **A**, then $|\mathbf{A}| = |\mathbf{L}| \cdot |\mathbf{U}|$ where $|\mathbf{L}| = l_{11}l_{22}l_{33}...l_{nn} \qquad |\mathbf{U}| = u_{11}u_{22}u_{33}...u_{nn}$

The number of arithmetic operations of LU decomposition is of order of *n*³. This is much less (in case of determinants of higher order) much less then *n*! operation necessary to perform if we use Lapalce expansion for calculation of determinant.

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Motivation

Many practical problems require to solve the large systems of linear equations **A x** = **b** , in which the matrix **A** is sparse, i.e. it has relatively small number of nonzero elements.

Standard elimination methods are not suitable for solving such large sparse linear systems.

(*Why*?)

Because during the elimination process we fill-up positions of originally zero elements – matrix is no more sparse. We chose an initial vector \mathbf{X}_0 and we generate a sequence of vectors

$$\mathbf{x}_0 \to \mathbf{x}_1 \to \mathbf{x}_2 \cdots,$$

which converge to the seeking solution \mathbf{X} .

Common feature of all iterative methods is fact, that each iteration step

$\mathbf{X}_k \rightarrow \mathbf{X}_{k+1}$

requires as many operations as multiplication of matrix **A** by a vector, which is for sparse matrices relatively small number of operations.

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Convergence

The most of classical iterative methods are based on decomposition of matrix $\mathbf{A} = \mathbf{M} - \mathbf{N}$, where **M** is regular matrix.

Then we can define a sequence of \mathbf{x}_{k} as

 $\mathbf{M}\mathbf{x}_{k+1} = \mathbf{N}\mathbf{x}_k + \mathbf{b},$

where initial approximation \mathbf{x}_0 is prescribed.

We say, that the **iterative method converge**, and we write $\mathbf{X}_k \to \mathbf{X}$, if numeral sequence $\|\mathbf{X}_k - \mathbf{X}\| \to 0.$

We denote $\mathbf{e}_k = \mathbf{X}_k - \mathbf{X}$ the error of k th iteration. Because $\mathbf{M}\mathbf{X} = \mathbf{N}\mathbf{X} + \mathbf{b}$, we get

$$\mathbf{M}(\mathbf{x}_{k+1} - \mathbf{x}) = \mathbf{N}(\mathbf{x}_k - \mathbf{x})$$
 or $\mathbf{e}_{k+1} = \mathbf{M}^{-1}\mathbf{N}\mathbf{e}_k$

If we denote $\mathbf{T} = \mathbf{M}^{-1}\mathbf{N}$ then we get $\|\mathbf{e}_{k+1}\| \leq \|\mathbf{T}\| \cdot \|\mathbf{e}_k\| \leq \|\mathbf{T}\|^2 \cdot \|\mathbf{e}_{k-1}\| \leq \cdots \leq \|\mathbf{T}\|^{k+1} \cdot \|\mathbf{e}_0\|$.

Convergence of iterative method is assured for any initial vector if $\|\mathbf{T}\| < 1$, where $\mathbf{T} = \mathbf{M}^{-1}\mathbf{N}$ is iteration matrix.

 $\|\mathbf{T}\|$ is an arbitrary matrix norm

"Entrywise" norm of matrix

$$\|\mathbf{A}\|_{p} = \|vec(\mathbf{A})\|_{p} = \left(\sum_{j=1}^{n} \sum_{j=1}^{n} |a_{ij}|^{p}\right)^{\frac{1}{p}}$$

$$\|\mathbf{A}\|_{F} = \sqrt{\sum_{i=1}^{n} \sum_{j=1}^{n} a_{ij}^{2}}$$

Frobenius norm (consistent with I_2)

$$\left\|\mathbf{A}\right\|_{\infty} = \max_{i,j} \left|a_{ij}\right|$$

max norm

Conditions for stopping iterations

How to decide, whether \mathbf{x}_{k+1} is good enough approximation of solution **x** ?

Usually we test one of the conditions:

1.
$$\|\mathbf{x}_{k+1} - \mathbf{x}_{k}\| \le \varepsilon \|\mathbf{x}_{k}\|$$

2. $\|\mathbf{r}_{k+1}\| \le \varepsilon (\|\mathbf{A}\| \cdot \|\mathbf{x}_{k+1}\| + \|\mathbf{b}\|)$

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Jacobi iterative method

Suppose that $\mathbf{A} = \mathbf{L} + \mathbf{D} + \mathbf{U}$, where \mathbf{D} is diagonal matrix, which has the same diagonal as \mathbf{A} , and where \mathbf{L} and \mathbf{U} are strictly lower and upper triangular parts of matrix \mathbf{A} , i.e.

$$\mathbf{D} = \begin{pmatrix} a_{11} & 0 & \cdots & 0 \\ 0 & a_{22} & & 0 \\ \vdots & & \ddots & \vdots \\ 0 & 0 & \cdots & a_{nn} \end{pmatrix}$$
$$\mathbf{L} = \begin{pmatrix} 0 & \cdots & 0 \\ a_{21} & 0 & & 0 \\ \vdots & \ddots & \ddots & \vdots \\ a_{n1} & \cdots & a_{n,n-1} & 0 \end{pmatrix}, \quad \mathbf{U} = \begin{pmatrix} 0 & a_{12} & \cdots & a_{1n} \\ \vdots & \ddots & \ddots & \vdots \\ 0 & & 0 & a_{n-1,n} \\ 0 & \cdots & \cdots & 0 \end{pmatrix}$$

Jacobi iterative method

Jacobi method is based on decomposition A = M - N, where M = D and N = -(L+U)and we write it as

$$\mathbf{D}\mathbf{x}_{k+1} = \mathbf{b} - (\mathbf{L} + \mathbf{U})\mathbf{x}_k$$

This system is easy to solve. If we write it in component form then

$$x_i^{(k+1)} = \frac{1}{a_{ii}} \left(b_i - \sum_{\substack{j=1\\j \neq i}}^n a_{ij} x_j^{(k)} \right) \ , \qquad i = 1, 2, ..., n \ .$$

Analysis of properties of iteration matrix $\mathbf{T} = -\mathbf{D}^{-1}(\mathbf{L} + \mathbf{U})$ leads to the statement, that Jacobi method converge, if A is diagonally dominant. Recall



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Gauss-Seidel method

Recall the Jacobi method

$$x_i^{(k+1)} = \frac{1}{a_{ii}} \left(b_i - \sum_{\substack{j=1\\j \neq i}}^n a_{ij} x_j^{(k)} \right), \qquad i = 1, 2, \dots, n$$

If we use
$$x_i^{(k+1)}$$
 instead of $x_i^{(k)}$

we get

Gauss-Seidel method

$$x_i^{(k+1)} = \frac{1}{a_{ii}} \left(b_i - \sum_{j=1}^{i-1} a_{ij} x_j^{(k+1)} - \sum_{j=i+1}^n a_{ij} x_j^{(k)} \right) , \qquad i = 1, 2, \dots, n .$$

In a matrix form

$$(\mathbf{D}+\mathbf{L})\mathbf{x}_{k+1} = \mathbf{b} - \mathbf{U}\mathbf{x}_k$$

Definition: Symmetric matrix ${f A}$ is positive-definite if

for any non-zero vector x it holds

 $x^T \cdot \mathbf{A} \cdot x > 0$

Analysis of properties of iteration matrix $\mathbf{T} = -(\mathbf{D} + \mathbf{L})^{-1} \mathbf{U}$ leads to the statement, that **Gauss-Seidel method converge**, **if A is diagonally dominant** or **positive-definite**. Checking, whether the matrix is positive-definite is usually problematic.

If we multiply any regular matrix **A** from left by their transpose matrix, the final matrix

$\mathbf{A}^T \mathbf{A}$

is symmetric and positive-definite.

Therefore, for system

$\mathbf{A}^T \mathbf{A} \mathbf{x} = \mathbf{A}^T \mathbf{b}$

it is assured that Gauss-Seidel method converge.

In such a case the convergence could be very slow.

- Convergence of Gauss-Seidel method is usually faster then convergence of Jacobi method
- There are matrices, for which the Gauss-Seidel method converge but Jacobi not and vice versa
- Jacobi method allows parallel processing, while the Gauss-Seidel method is sequential from their core

- Gauss elimination method and Cramer's rule lead to the solution. (Without round off errors we could find the exact solution.)
- The base of GEM is the modification of matrix to triangular form. (using elementary row operations)
- The influence of round-off errors on GEM could be considerable, therefore we use partial pivoting.
- The GEM is demanding from time and memory aspects. It is best suited for not very large systems with dense matrix.
- The Cramer's rule is suitable only for very small systems.

- Using the iterative method we usually find only approximate solution.
- At the beginning we chose the initial approximation of solution and we refine the solution by repeatedly inserting it into iteration formula.

The computation is usually finished when the norm of difference of two consecutive iterations is small enough.

- Iterative methods could also diverge. This depend on the properties of the matrix.
- Iterative methods are suitable for solving large systems with sparse matrix.
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Jacobi and Gauss-Seidel methods are so called stationary iterative methods.

Stationary iterative methods solve a linear system with an operator approximating the original one; and based on a measurement of the error in the result, form a "correction equation" for which this process is repeated.

While these methods are simple to derive, implement, and analyze, convergence is only guaranteed for a limited class of matrices.

Linear stationary iterative methods are also called **relaxation methods**.

Krylov subspace methods

work by forming a basis of the sequence of successive matrix powers times the initial residual (the **Krylov sequence**).

The approximations to the solution are then formed by minimizing the residual over the subspace formed. The prototypical method in this class is the **conjugate gradient method** (CG).

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The **conjugate gradient method** is implemented as an iterative algorithm, applicable to large sparse systems with symmetric, positive-definite matrix.

> We say that two non-zero vectors **u** and **v** are conjugate (with respect to **A**) if

$$\mathbf{u}^{\mathrm{T}}\mathbf{A}\mathbf{v}=0$$

The iterative algorithms is described by this formulas:

$$\begin{split} \mathbf{v}^{(1)} &= \mathbf{r}^{(1)} = \mathbf{b} - \mathbf{A} \mathbf{x}^{(1)} \ , \\ \alpha_k &= \frac{\mathbf{v}^{(k)T} \mathbf{r}^{(k)}}{\mathbf{v}^{(k)T} \mathbf{A} \mathbf{v}^{(k)}} \ , \\ \mathbf{x}^{(k+1)} &= \mathbf{x}^{(k)} + \alpha_k \mathbf{v}^{(k)} \ , \\ \mathbf{r}^{(k+1)} &= \mathbf{r}^{(k)} - \alpha_k \mathbf{A} \mathbf{v}^{(k)} \ , \\ \beta_k &= -\frac{\mathbf{v}^{(k)T} \mathbf{A} \mathbf{r}^{(k+1)}}{\mathbf{v}^{(k)T} \mathbf{A} \mathbf{v}^{(k)}} \ , \\ \mathbf{v}^{(k+1)} &= \mathbf{r}^{(k+1)} + \beta_k \mathbf{v}^{(k)} \ . \end{split}$$

Conjugate gradient method (CG)

From the fundamentals of algorithm it follows that after *n* iterations we obtain exact solution of the system and therefore **it is not** an iterative method in a strict sense.

This would be true only if there are no round-off errors.

Therefore we have to look at CG method as an iterative method and we have to define stop criteria.

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Generalizations of CG method

In case of nonsymmetrical and not necessary positive-definite matrices we use **biconjugate gradient method** (e.g. subroutine linbcg from Numerical Recipes).

The ordinary conjugate gradient method is their special case.

The other variant suitable for symmetric but not positive-definite matrices \mathbf{A} we obtain using exchange of all matrix multiplications $\mathbf{a} \cdot \mathbf{b}$ for $\mathbf{a} \cdot \mathbf{A} \cdot \mathbf{b}$. This method is called **algorithm of minimal residual**, because we minimize the form $Q(\mathbf{x}) = \frac{1}{2}\mathbf{r}^T \cdot \mathbf{r} = \frac{1}{2}|\mathbf{A} \cdot \mathbf{x} \cdot \mathbf{b}|^2$.

Generalization for nonsymmetrical matrices leads to the **generalized minimal residual method (GMRES)**.

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Let $\mathbf{x}^{(k)}$ be an approximate solution in the k-th step of CG and let \mathbf{x}^* be the exact solution.

For symmetric, positive-definite matrix we define **condition number of matrix** κ(A) and A-norm of an arbitrary vector z:

$$\kappa(\mathbf{A}) \coloneqq \frac{\lambda_{\max}(\mathbf{A})}{\lambda_{\min}(\mathbf{A})}, \qquad \|\mathbf{z}\|_{\mathbf{A}} \coloneqq (\mathbf{A}\mathbf{z}, \mathbf{z})^{1/2}$$

$$\begin{split} & \left\| \mathbf{x}^{*} - \mathbf{x}^{(k)} \right\|_{\mathbf{A}} \leq 2 \! \left[\frac{\sqrt{\kappa(\mathbf{A})} - 1}{\sqrt{\kappa(\mathbf{A})} + 1} \right]^{k} \left\| \mathbf{x}^{*} - \mathbf{x}^{(1)} \right\|_{\mathbf{A}} \end{split}$$

If $\kappa(\mathbf{A}) >> 1$ the convergence is **very slow**.

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The aim of preconditioning is to speed-up the convergence of iterative method in such a way, that we solve an alternative system of linear equations in which the coefficient matrix has

lower condition number $\boldsymbol{\mathcal{K}}$

then the original coefficient matrix.

Let $\mathbf{A} \cdot \mathbf{x} = \mathbf{b}$ be the original linear system.

Left preconditioning is defined as: Let matrix M is regular and "close" to matrix A. Then we solve the system $M^{-1}Ax = M^{-1}b$

Right preconditioning is defined as: Let matrix M is regular. Then we solve system $AM^{\text{-1}}u=b, \ \ x=M^{-1}u \ .$

Preconditioning in CG method

Original algorithms CG:

 $\mathbf{r}^{(1)} = \mathbf{b} - \mathbf{A}\mathbf{x}^{(1)}$ $\mathbf{v}^{(1)} = \mathbf{r}^{(1)}$

Algorithm of preconditioned CG:

$$\mathbf{r}^{(1)} = \mathbf{b} - \mathbf{A}\mathbf{x}^{(1)}$$
, $\mathbf{z}^{(1)} = \mathbf{M}^{-1}\mathbf{r}^{(1)}$
 $\mathbf{v}^{(1)} = \mathbf{z}^{(1)}$

$$\alpha_k = \frac{\mathbf{v}^{(k)T} \mathbf{r}^{(k)}}{\mathbf{v}^{(k)T} \mathbf{A} \mathbf{v}^{(k)}}$$

$$\mathbf{x}^{(k+1)} = \mathbf{x}^{(k)} + \alpha_k \mathbf{v}^{(k)}$$

$$\mathbf{r}^{(k+1)} = \mathbf{r}^{(k)} - \alpha_k \mathbf{A} \mathbf{v}^{(k)}$$

$$\alpha_k = \frac{\mathbf{z}^{(k)T} \mathbf{r}^{(k)}}{\mathbf{v}^{(k)T} \mathbf{A} \mathbf{v}^{(k)}}$$
$$\mathbf{x}^{(k+1)} = \mathbf{x}^{(k)} + \alpha_k \mathbf{v}^{(k)}$$

$$\mathbf{r}^{(k+1)} = \mathbf{r}^{(k)} - \alpha_k \mathbf{A} \mathbf{v}^{(k)}$$
, $\mathbf{z}^{(k+1)} = \mathbf{M}^{-1} \mathbf{r}^{(k+1)}$

$$\beta_k = -\frac{\mathbf{v}^{(k)T} \mathbf{A} \mathbf{r}^{(k+1)}}{\mathbf{v}^{(k)T} \mathbf{A} \mathbf{v}^{(k)}} \qquad \beta_k = \frac{\mathbf{z}^{(k+1)T} \mathbf{r}^{(k+1)}}{\mathbf{z}^{(k)T} \mathbf{r}^{(k)}}$$

$$\mathbf{v}^{(k+1)} = \mathbf{r}^{(k+1)} + \beta_k \mathbf{v}^{(k)} \qquad \mathbf{v}^{(k+1)} = \mathbf{z}^{(k+1)} + \beta_k \mathbf{v}^{(k+1)}$$

Jacobi preconditioner

One of the simplest forms of preconditioning, is obtained by the choosing the preconditioner to be the diagonal of the matrix matice ${\bf A}$

$$M_{ij} = \begin{cases} A_{ii} & \text{if } i = j \\ 0 & \text{if } i \neq j \end{cases}$$

Then

$$M_{ij}^{-1} = \frac{\delta_{ij}}{A_{ii}}$$

This preconditioning we call Jacobi preconditioner or

diagonal scaling.

Advantages of Jacobi preconditioner are the easy implementation and low memory requirements. Other types of preconditioning

More sophisticated choices of preconditioner have to be a compromise between the reduction of condition number of the matrix and therefore faster convergence and

the time necessary to calculate inverse matrix ${f M}^{-1}$.

More information here

M. Benzi (2002): **Preconditioning Techniques for Large Linear Systems: A Survey** *Journal of Computational Physics* **182**, 418-477 <u>doi: 10.1006/jcph.2002.7176</u>

Example



Example



Figure: PCG, $\kappa(A) = 8 * 10^9$.

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 - SVD for more equations than unknowns

Lets have a square matrix \mathbf{A} . Its **eigenvalues** will be denoted as λ_n , right and left eigenvectors as x and y, and it holds $\det(\mathbf{A} - \lambda \mathbf{I}) = 0$ $\mathbf{A}\mathbf{x}_i = \lambda_i \mathbf{x}_i$ $\mathbf{y}_i^T \mathbf{A} = \lambda_i \mathbf{y}_i^T$ If $\mathbf{A} \in \Re$ and is symmetric, then $\lambda_{i}, \mathbf{x}_{i} \equiv \mathbf{y}_{i} \in \Re$ and eigenvectors generate the orthogonal base. Affine transform $\mathbf{P}^{-1}\mathbf{A}\mathbf{P}$

does not change the eigenvalues of matrix \mathbf{A} .

Real **orthogonal matrix** is such a square matrix \mathbf{Q} , for which its transpose is equal to its inverse.

 $\mathbf{Q}^T \cdot \mathbf{Q} = \mathbf{Q} \cdot \mathbf{Q}^T = \mathbf{I}$

If ${f Q}$ is not a square matrix, then conditions

$$\mathbf{Q}^T \cdot \mathbf{Q} = \mathbf{I}$$
 and $\mathbf{Q} \cdot \mathbf{Q}^T = \mathbf{I}$

are not equivalent.

Condition $\mathbf{Q}^T \cdot \mathbf{Q} = \mathbf{I}$ says that matrix \mathbf{Q} is matrix with orthogonal columns.

Condition $\mathbf{Q} \cdot \mathbf{Q}^T = \mathbf{I}$ says that matrix \mathbf{Q} is matrix with orthogonal rows.

If the number of equations M is less than the number of unknowns N, or if M = N but equations are linearly dependent, then the system has no solution or it has more than one solution.

In the later case the space of solutions is given by a particular solution added to the any linear combination of N – M vectors.

The task to find the space of solution for matrix A is possible to solve using **singular value decomposition of matrix A**.

If the number of equations M is greater than the number of unknowns N, then in general there is no solution vector and the system is called **overdetermined**.

However, we could find a best "compromise" solution, which is "the closest" solution that satisfy all equations.

If "the closest" we define in a sense of least square, i.e. the sum of square of residuals is the smallest possible, then the overdetermined system is reduced to solvable problem called the **least square method**. Reduced system of equations could be written as system $N \ge N$ equation $(\mathbf{A}^T \cdot \mathbf{A}) \cdot \mathbf{x} = (\mathbf{A}^T \cdot \mathbf{b})$.

This equations we call **normal equations** of a least square problem.

Singular value decomposition has many common features with the least square method, which we show later.

Direct solution of normal equations is in general not the best way to find a solution of least square problem.

Singular value decomposition

In many cases, when GEM or LU decomposition fail, singular value decomposition (SVD) precisely diagnose, where is the problem and in many cases it also offer a suitable numerical solution.

SVD is also method to solve many least square problems.

SVD is based on the following theorem of linear algebra:

Each matrix ${f A}$ of type M x N,

which the number of rows M is greater or equal to the number of columns N, could be decomposed to a product of

matrix with orthogonal columns ${f U}$ of type M x N,

diagonal matrix ${f W}$ of type N x N

with positive or zero entries (singular values)

and transpose orthogonal matrix \mathbf{V} of type $N \ge N$.





SVD could be done also if M < N.

In such a case the singular values w_j for j = M+1, ..., Nare all zero

as well as corresponding columns of matrix ${f U}_{\cdot}$

There is a lot of algorithms of SVD, proven is subroutine **svdcmp** from *Numerical Recipes*.

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Singular value decomposition of a square matrix

If ${f A}$ is a square matrix of type N x $N_{
m c}$

then \mathbf{U} , \mathbf{W} and \mathbf{V} are all square matrices of type $N \ge N$.

Because \mathbf{U} and \mathbf{V} are orthogonal, their inverse are equal to transpose.

Then we can write a formula for inverse matrix \mathbf{A} $\mathbf{A}^{-1} = \mathbf{V} \cdot \left[\operatorname{diag} \left(1 / w_j \right) \right] \cdot \mathbf{U}^T$

> SVD offers clear diagnosis of situation, if some of singular values are zero or close to zero.

Recall one of definition of condition number of matrix $\kappa(A)$:

$$\kappa(\mathbf{A}) \coloneqq \|\mathbf{A}\| \|\mathbf{A}^{-1}\| = \frac{\max\left\{w_j\right\}}{\min\left\{w_j\right\}}$$

The matrix is **singular** if its condition umber is infinity.

The matrix is **ill-conditioned** if reciprocal value of its condition number is close to the machine epsilon of a computer, i.e. less than 10⁻⁶ for single precision or 10⁻¹² for double precision. The case of system $\mathbf{A} \cdot \mathbf{x} = \mathbf{b} \ ,$

in which the coefficient matrix ${f A}$ is singular:

At first, lets have a look at the homogeneous case, i.e. the case of b=0. In other words

we are looking for null space of \boldsymbol{A}

$$null(\mathbf{A}) = \left\{ \mathbf{x} \in \mathbb{R}^{n} : \mathbf{A}\mathbf{x} = \mathbf{0} \right\} = \left\{ \mathbf{x} \in \mathbb{R}^{n} : \mathbf{U}\mathbf{w}\mathbf{V}^{T}\mathbf{x} = \mathbf{0} \right\}$$
$$= \left\{ \mathbf{x} \in \mathbb{R}^{n} : \mathbf{U}^{T}\mathbf{U}\mathbf{w}\mathbf{V}^{T}\mathbf{x} = \mathbf{0} \right\} = \left\{ \mathbf{x} \in \mathbb{R}^{n} : \mathbf{w}\mathbf{V}^{T}\mathbf{x} = \mathbf{0} \right\}$$

The case of system $\mathbf{A} \cdot \mathbf{x} = \mathbf{b}$,

in which the coefficient matrix ${f A}$ is singular:

At first, lets have a look at the homogeneous case, i.e. the case of b=0. In other words we are looking for null space of A

SVD gives direct solution – each column of matrix ${f V}$,

which corresponding singular value w_j is zero is a solution.

The case of system $\mathbf{A} \cdot \mathbf{x} = \mathbf{b} \ ,$

in which the coefficient matrix ${f A}$ is singular:

Now lets have a look at range of matrix ${f A}$

$$range(\mathbf{A}) = \left\{ \mathbf{A}\mathbf{x} : \mathbf{x} \in \mathbb{R}^n \right\} = \left\{ \mathbf{U}\mathbf{w}\mathbf{V}^T\mathbf{x} : \mathbf{x} \in \mathbb{R}^n \right\} = \left\{ \mathbf{U}\mathbf{w}\mathbf{y} : \mathbf{y} \in \mathbb{R}^n \right\}$$
$$\mathbf{y} = \mathbf{V}^T\mathbf{x}$$

```
The case of system \mathbf{A} \cdot \mathbf{x} = \mathbf{b},
```

in which the coefficient matrix ${f A}$ is singular:

Now lets have a look at range of matrix ${f A}$

The range of matrix A is composed by span of columns of matrix U, which corresponding singular value w_j is nonzero.
Singular value decomposition of a square matrix

The solution of a system with nonzero right-hand-side using SVD is this:

- we exchange $1/w_j$ by zero if $w_j=0$
- then we calculate (from right to left)

$$\mathbf{x} = \mathbf{V} \cdot \left[\operatorname{diag} \left(1 / w_j \right) \right] \cdot \left(\mathbf{U}^T \cdot \mathbf{b} \right)$$

If a particular solution lies in range of \mathbf{A} , then it has the smallest size $|\mathbf{x}|^2$.

If a particular solution does not lies in range of ${f A}$,

then \mathbf{x} minimize the residuum of solution $r \coloneqq |\mathbf{A} \cdot \mathbf{x} - \mathbf{b}|$.

Singular value decomposition of a square matrix

Matrix \mathbf{A} is not singular



Singular value decomposition of a square matrix



Up till now we considered only extreme cases that the coefficient matrix is or is not singular.

There is often the case

that the singular values w_j are very small but nonzero, so the matrix is ill-conditioned.

In such a case direct methods can offer formally the solution, but the solution vector has unreasonably large entries, which during the algebraic manipulations with matrix ${\bf A}$ leads to very bad approximation of the right-hand-side vector.

At that time is better small values w_j set to zero and the solution calculate using (with replace of $1/w_j$ by zero if $w_j=0$) $\mathbf{x} = \mathbf{V} \cdot \left[\operatorname{diag} \left(1/w_j \right) \right] \cdot \left(\mathbf{U}^T \cdot \mathbf{b} \right)$

We have to be cautious and we have to chose the good threshold level for zeroise of W_j

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If we have less equation than unknowns we expect N - M dimensional space of solutions.

In this case SVD offers N - M zero or negligible small values of w_j . If some of M equations degenerate, then we could have additional zero-valued w_j .

Then those columns of ${f V}$,

which corresponds to zero singular value W_j makes a basis vectors of seeking solution space.

Particular solution could by find using

$$\mathbf{x} = \mathbf{V} \cdot \left[\operatorname{diag} \left(1 / w_j \right) \right] \cdot \left(\mathbf{U}^T \cdot \mathbf{b} \right)$$

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Singular value decomposition for more equations than unknowns

If we have more equations than unknowns we are seeking solution in a least square sense.

We are solving system written as:





In this case usually it is not necessary set to zero values w_j however the unusually small values indicate that the data are not sensitive to some parameters.

Lecture 4

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- 2. Iterative method for solving linear systems
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- 4. Jacobi method, Gauss-Seidel method
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